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To cite this article: Medha Itagi Huilgol , M. Rajeshwari & S. Syed Asif Ulla (2012) Products of distance degree regular and distance degree injective graphs, Journal of Discrete Mathematical Sciences and Cryptography, 15:4-5, 303-314, DOI: [10.1080/09720529.2012.10698382](https://doi.org/10.1080/09720529.2012.10698382)

To link to this article: <https://doi.org/10.1080/09720529.2012.10698382>



Published online: 03 Jun 2013.



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## Products of distance degree regular and distance degree injective graphs

Medha Itagi Huilgol<sup>1,\*</sup>

M. Rajeshwari<sup>2,†</sup>

S. Syed Asif Ulla<sup>1,§</sup>

<sup>1</sup>*Department of Mathematics  
Bangalore University  
Central College Campus  
Bangalore 560 001  
India*

<sup>2</sup>*Department of Mathematics  
Acharya Institute of Technology  
Bangalore 560 090  
India*

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### Abstract

The eccentricity  $e(u)$  of a vertex  $u$  is the maximum distance of  $u$  to any other vertex in  $G$ . The distance degree sequence (dds) of a vertex  $v$  in a graph  $G = (V, E)$  is a list of the number of vertices at distance  $1, 2, \dots, e(u)$  in that order, where  $e(u)$  denotes the eccentricity of  $u$  in  $G$ . Thus the sequence  $(d_{i_0}, d_{i_1}, d_{i_2}, \dots, d_{i_j}, \dots)$  is the dds of the vertex  $v_i$  in  $G$  where  $d_{i_j}$  denotes number of vertices at distance  $j$  from  $v_i$ . A graph is distance degree regular (DDR) graph if all vertices have the same dds. A graph is distance degree injective (DDI) graph if no two vertices have same dds.

In this paper we consider Cartesian and normal products of DDR and DDI graphs. Some structural results have been obtained along with some characterizations.

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**Keywords:** *Distance degree sequence, Distance degree regular (DDR) graphs, Distance degree injective (DDI) graphs, Cartesian and Normal product of graphs.*

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\*E-mail: medha@bub.ernet.in

†E-mail: rajeshwarim@acharya.ac.in

§E-mail: syedasif.ulla84@gmail.com

## 1. Introduction

Unless mentioned otherwise for terminology and notation the reader may refer Buckley and Harary [4], new ones will be introduced as and when found necessary.

Products in graphs have always given more generalized results compared to the graphs involved in the product itself. It is also a powerful tool to construct bigger graphs given smaller ordered (sized) graphs. Many parameters are tested in the products in literature [6], [7], [8], etc. Among many products defined between graphs the cartesian product is the most used one. Recently a whole monograph by Imrich *et. al.*, [10] is dedicated to graphs and their cartesian product. The cartesian product is defined as,

The cartesian product of two graphs  $G$  and  $H$ , denoted  $G \square H$ , is a graph with vertex set  $V(G \square H) = V(G) \times V(H)$ , that is, the set  $g \in V(G) \& h \in V(H)$ .

The edge set of  $G \square H$  consists of all pairs  $[(g_1, h_1), (g_2, h_2)]$  of vertices with  $[g_1, g_2] \in E(G)$  and  $h_1 = h_2$  or  $g_1 = g_2$  and  $[h_1, h_2] \in E(H)$ .

Also the normal product is defined as,

The normal product of two graphs  $G$  and  $H$ , denoted  $G \oplus H$ , is a graph with vertex set  $V(G \oplus H) = V(G) \times V(H)$ , that is, the set  $g \in V(G), h \in V(H)$  and an edge  $[(g_1, h_1), (g_2, h_2)]$  exists whenever any of the following conditions hold good :

- (i)  $[g_1, g_2] \in E(G)$  and  $h_1 = h_2$ ,
- (ii)  $g_1 = g_2$  and  $[h_1, h_2] \in E(H)$ ,
- (iii)  $[g_1, g_2] \in E(G) \& [h_1, h_2] \in E(H)$ .

The distance  $d(u, v)$  from a vertex  $u$  of  $G$  to a vertex  $v$  is the length of a shortest  $u$  to  $v$  path. The eccentricity  $e(v)$  of  $v$  is the distance to a farthest vertex from  $v$ . If,  $dist(u, v) = e(u), (v \neq u)$ , we say that  $v$  is an eccentric vertex of  $u$ . The radius  $r(G)$  is the minimum eccentricity of the vertices, whereas the diameter  $d(G)$  is the maximum eccentricity. A vertex  $v$  is a central vertex if  $e(v) = r(G)$ , and a vertex is an antipodal vertex if  $e(v) = d(G)$ . A graph is self-centered if every vertex has the same eccentricity, i.e.,  $r(G) = d(G)$ .

The distance degree sequence (*dds*) of a vertex  $v$  in a graph  $G = (V, E)$  is a list of the number of vertices at distance  $1, 2, \dots, e(v)$  in that order, where  $e(v)$  denotes the eccentricity of  $v$  in  $G$ . Thus, the sequence  $(d_{i_0}, d_{i_1}, d_{i_2}, \dots, d_{i_j}, \dots)$  is the dds of the vertex  $v_i$  in  $G$  where,  $d_{i_j}$  denotes

number of vertices at distance  $j$  from  $v_i$ . The concept of distance degree regular (DDR) graphs was introduced by G. S. Bloom *et. al.*, [1], as the graphs for which all vertices have the same dds. For example, the three dimensional cube  $Q_3 = K_2 \times K_2 \times K_2$  is a DDR graph with each vertex having its dds as  $(1,3,3,1)$ . Clearly,  $d_{i_1}$  denotes the degree of the vertex  $v_i$  in  $G$  and hence, in general, a DDR graph must be a regular graph; but, it is easy to verify that a regular graph may not be DDR. In Bloom [1], detailed study of DDR graphs can be found and one of the fundamental results therein states that "Every regular graph with diameter at most two is DDR". Bloom, Quintas and Kennedy [2] have dealt many problems concerning distance and path degree sequences in graphs. Halberstam *et. al.*, [9] have dealt in particular the distance and path degree sequences for cubic graphs. It is worth to mention that computer investigation and generation of cubic graphs is done by Brinkmann [3] and Bussemaker *et. al.*, [5]. In [12], Itagi Huilgol *et. al.*, have listed all DDR graphs of diameter three with extremal degree regularity. In the same paper they have shown the existence of a diameter three DDR graph of any arbitrary degree regularity. In [13], Itagi Huilgol *et. al.*, have constructed DDR graphs of arbitrary diameter. Also, they have studied the DDR graphs with respect to other parameters.

A graph is distance degree injective (DDI) graph if no two vertices have same dds. These graphs were defined by G.S. Bloom *et. al.*, in [2]. DDI graphs being highly irregular, in comparison with the DDR graphs, at least the degree regularity is looked into by Jiri volf in [11]. A particular case of cubic DDI graphs is considered by Martenez and Quintas in [14]. There are very few examples of DDI graphs, so it is important to get DDI graphs from smaller (sized/ordered) DDI or other graphs as products.

In this paper we consider cartesian and normal products of DDR and DDI graphs. For some products both necessary and sufficient conditions have been obtained.

## 2. Cartesian product of DDR and DDI graph

In this section we consider cartesian products of DDR and DDI graphs.

**Theorem 2.1:** *Cartesian product of two graphs  $G_1$  and  $G_2$  is a DDR graph if and only if both  $G_1$  and  $G_2$  are DDR graphs.*

**Proof.** Let  $G_1$  and  $G_2$  be two DDR graphs having the dds of each vertex  $(d_0, d_1, d_2, \dots, d_{r_1})$  and  $(d'_0, d'_1, d'_2, \dots, d'_{r_2})$  respectively, where  $r_1$  and  $r_2$  are radii of  $G_1$  and  $G_2$  respectively. In the cartesian product of any two graphs, the distance between any two vertices  $(u_1, v_1)$  and  $(u_2, v_2)$  is given by  $d_{G_1 \square G_2}((u_1, v_1), (u_2, v_2)) = d_{G_1}(u_1, u_2) + d_{G_2}(v_1, v_2)$  as in [15]. Now let  $u$  be any vertex in  $G_1$  and  $v$  be any vertex in  $G_2$ . Then, it is very clear that the number of vertices at distance  $i$  from  $(u, v)$  in  $G_1 \square G_2$

$$d_{i_{G_1 \square G_2}}(u, v) = d_i(u) + d'_i(v) + \sum_{j=1}^{i-1} d_j(u) d'_{i-j}(v).$$

Since the graphs  $G_1$  and  $G_2$  are DDR graphs  $d_i(u) = d_i(x)$ ,  $0 \leq i \leq \text{diam}(G_1)$  for all  $x \in G_1$  and  $d'_i(v) = d'_i(y)$ ,  $0 \leq i \leq \text{diam}(G_2)$  for all  $y \in G_2$ ,  $d_{i_{G_1 \square G_2}}(u, v) = d_{i_{G_1 \square G_2}}(s, t)$ ,  $0 \leq i \leq \text{diam}(G_1) + \text{diam}(G_2)$  for all  $(s, t) \in G_1 \square G_2$ . Hence the graph  $G_1 \square G_2$  is a DDR graph.

Now let  $G_1 \square G_2$ , the cartesian product of  $G_1$  and  $G_2$  be a DDR graph. Suppose  $G_1$  is not a DDR graph, then there exist at least two vertices  $u$  and  $v$  having different dds i.e.,  $(d_0(u), d_1(u), d_2(u), \dots, d_{e(u)}(u))$  and  $(d_0(v), d_1(v), d_2(v), \dots, d_{e(v)}(v))$  are dds of  $u$  and  $v$ , respectively in  $G_1$  and  $k$ , the minimum value of  $i$ ,  $1 \leq i \leq d(G_1)$ , such that  $d_k(u) \neq d_k(v)$ . Let  $w$  be any vertex in  $G_2$ , having the dds  $(d'_0(w), d'_1(w), d'_2(w), \dots, d'_{e(w)}(w))$  and  $d'_k(w)$  be the number of vertices at distance  $k$  from  $w$  in  $G_2$ . The number of vertices at distance  $k$  from  $(u, w)$  in  $G_1 \square G_2$  is given by  $d_{k_{G_1 \square G_2}}(u, w) = d_k(u) + d'_k(w) + d_1(u) d'_{(k-1)}(w) + d_2(u) d'_{(k-2)}(w) + d_3(u) d'_{(k-3)}(w) + \dots + d_{(k-1)}(u) d'_1(w)$  and the number of vertices at distance  $k$  from  $(v, w)$  in  $G_1 \square G_2$  is given by  $d_{k_{G_1 \square G_2}}(v, w) = d_k(v) + d'_k(w) + d_1(v) d'_{(k-1)}(w) + d_2(v) d'_{(k-2)}(w) + d_3(v) d'_{(k-3)}(w) + \dots + d_{(k-1)}(v) d'_1(w)$ . Hence  $d_{k_{G_1 \square G_2}}(u, w) \neq d_{k_{G_1 \square G_2}}(v, w)$ , since  $d_k(u) \neq d_k(v)$  and  $d_j(u) = d_j(v)$ , for all  $j$ ,  $0 \leq j \leq k$ . Hence  $G_1 \square G_2$  is non-DDR graph, a contradiction. Hence  $G_1$  should be a DDR graph. Similarly we can prove  $G_2$  is also a DDR graph. Hence, the result.  $\square$

**Theorem 2.2:** *If the cartesian product of two graphs  $G_1$  and  $G_2$  is DDI then both  $G_1$  and  $G_2$  are DDI graphs.*

**Proof.** Let  $G_1 \square G_2$  be a DDI graph. Suppose  $G_1$  is not a DDI graph, then there exist at least two vertices  $u_1, u_2$  in  $G_1$  having the same dds, i.e.,  $dds(u_1) = dds(u_2)$ . Let  $v_1$  be any vertex in  $G_2$ . Then the number of vertices at distance  $l$ ,  $0 \leq l \leq e(u_1) + e(v_1)$  from  $(u_1, v_1)$  is given by

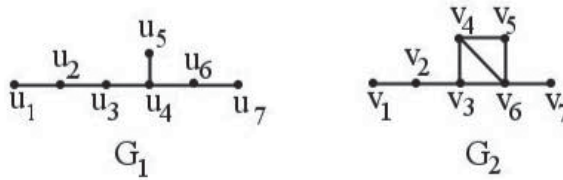
$$d_{l_{G_1 \square G_2}}(u_1, v_1) = d_i(u_1) + d'_i(v_1) + \sum_{j=1}^{i-1} d_j(u_1) d'_{i-j}(v_1).$$

and the number of vertices at distance  $l, 0 \leq l \leq e(u_2) + e(v_1)$  from  $(u_2, v_1)$  is given by

$$d_{l_{G_1 \square G_2}}(u_2, v_1) = d_i(u_2) + d'_i(v_1) + \sum_{j=1}^{i-1} d_j(u_2) d'_{i-j}(v_1).$$

Since,  $dds(u_1) = dds(u_2)$ , we get  $d_{l_{G_1 \square G_2}}(u_1, v_1) = d_{l_{G_1 \square G_2}}(u_2, v_1)$ , for all  $l, 0 \leq l \leq e(u_2) + e(v_1)$ , i.e.,  $dds(u_1, v_1) = dds(u_2, v_1)$ , hence  $G_1 \square G_2$  is not DDI, a contradiction. Hence  $G_1$  should be DDI. Similarly, we can prove  $G_2$  is also DDI. Hence, the proof.  $\square$

**Remark 1.** Cartesian product of two DDI graphs need not be DDI. The following are the two DDI graphs whose cartesian product is not DDI.



**Figure 1**  
Two DDI graphs whose cartesian product is not DDI

**Lemma 2.1:** Let  $G_1$  and  $G_2$  be two DDI graphs. Let  $A = \{dds(u_i) \mid u_i \in V(G_1)\}$  and  $B = \{dds(v_i) \mid v_i \in V(G_2)\}$ . If  $|A \cap B| \geq 2$  then  $G_1 \square G_2$  is not DDI.

**Proof.** Let  $|A \cap B| \geq 2$  Then there exist  $u_k, u_l$  in  $G_1$  and  $v_m, v_n$  in  $G_2$  such that  $dds(u_k) = dds(v_m)$  and  $dds(u_l) = dds(v_n)$ . Hence in  $G_1 \square G_2$ ,  $dds(u_k, v_n) = dds(u_l, v_m)$ , making  $G_1 \square G_2$  non DDI. Hence, the proof  $\square$

**Theorem 2.3.** Let  $G_1$  and  $G_2$  be any two graphs. Let  $u$  be any vertex in  $G_1$  and  $S$  be a subset of  $V(G_2)$  such that no two vertices of  $S$  have same  $dds$ , then no two vertices of  $\{u\} \times S$  have same  $dds$  in  $G_1 \square G_2$ .

**Proof.** Let  $u$  be any vertex in  $G_1$  and  $S$  be a subset of  $V(G_2)$  such that no two vertices of  $S$  have same  $dds$ . Suppose there exist at least two vertices  $(u, v)$  and  $(u, w)$  in  $\{u\} \times S$  having same  $dds$ . Hence  $d_{l_{G_1 \square G_2}}(u, v) = d_{l_{G_1 \square G_2}}(u, w)$ , for all  $l, 0 \leq l \leq e(u) + e(v)$ , here  $e(v) = e(w)$ .

$$\text{Hence } d_l(u) + d'_l(v) + \sum_{j=1}^{l-1} d_j(u) d'_{l-j}(v) = d_l(u) + d'_l(w) + \sum_{j=1}^{l-1} d_j(u) d'_{l-j}(w),$$

$$\text{implies } d'_l(v) + \sum_{j=1}^{l-1} d_j(u) d'_{l-j}(v) = d'_l(w) + \sum_{j=1}^{l-1} d_j(u) d'_{l-j}(w) \rightarrow (1)$$

For  $l = 1$ , eq.(1) implies  $d'_1(v) = d'_1(w)$ .

For  $l = 2$ , eq.(1) implies  $d'_2(v) + d_1(u) d'_1(v) = d'_2(w) + d_1(u) d'_1(w)$

$\Rightarrow d'_2(v) = d'_2(w)$  and so on

For,  $l = e(v) = e(w)$ , eq.(1) implies  $d'_{e(v)}(v) = d'_{e(w)}(w)$ . Hence,  $dds(v) = dds(w)$ , a contradiction. Hence no two vertices of  $\{u\} \times S$  have same  $dds$  in  $G_1 \square G_2$ . Hence, the proof.  $\square$

**Remark 2.** Let  $S_1$  and  $S_2$  be two subsets of  $V(G_1)$  such that every pair  $(x, y)$ ,  $x \in S_1, y \in S_2$  satisfies  $dds(x) \neq dds(y)$  and let  $z \in V(G_2)$  be any vertex in  $G_2$ , then in  $G_1 \square G_2$ , the subsets  $\{(x, z) \mid x \in S_1\}$  and  $\{(y, z) \mid y \in S_2\}$  are such that every pair  $((x, z), (y, z))$ ,  $x \in S_1$  and  $y \in S_2$  satisfies  $dds(x, z) \neq dds(y, z)$ .

### 3. Normal product of DDR and DDI graphs

In this section we consider normal product of DDR and DDI graphs. Stevanović in [16] has considered the distance between any pair of vertices in normal product. Given two vertices  $(u_i, v_j)$  and  $(u_k, v_m)$  the distance between these two vertices in the normal product is given by;

$$d_{G_1 \oplus G_2}((u_i, v_j), (u_k, v_m)) = \max\{d_{G_1}(u_i, v_k), d_{G_2}(u_j, v_m)\}$$

Immediate conclusion we can draw as follows:

**Lemma 3.1:** Let  $S = \{G_i \mid i \geq 2\}$ . If there exist  $k$  such that  $G_k$  is self centred and  $\text{diam}(G_k) \geq \text{diam}(G_i)$  for all  $i \geq 2$  then normal product of all the graphs in  $S$  is a self centred graph with diameter equal to  $\text{diam}(G_k)$ .

**Theorem 3.1:** Normal product  $G_1 \oplus G_2$  of two graphs  $G_1$  and  $G_2$  is DDR if and only if both  $G_1$  and  $G_2$  are DDR graphs.

*Proof.* Let  $G_1$  and  $G_2$  be two DDR graphs having the dds  $(d_0, d_1, d_2, \dots, d_{d(G_1)})$  and  $(d'_0, d'_1, d'_2, \dots, d_{d(G_2)})$  respectively, then the number of vertices at distance  $i, 0 \leq i \leq \max \{diam(G_1), diam(G_2)\}$  from any vertex  $u_l, v_m$  in  $G_1 \oplus G_2$  is given by

$$d_{i_{G_1 \oplus G_2}}(u_l, v_m) = d_i(u_l) \cdot d'_i(v_m) + d_i(u_l) \sum_{j=0}^{i-1} d'_j(v_m) + d'_i(v_m) \sum_{j=0}^{i-1} d_j(u_l).$$

Hence, the Normal product  $G_1 \oplus G_2$  is DDR.

Conversely, let  $G_1 \oplus G_2$  be DDR. Suppose  $G_1$  is not DDR, then there exist at least two vertices  $u_1$  and  $u_2$  in  $G_1$  such that  $dds(u_1) \neq dds(u_2)$ . Let  $k$  be the minimum value such that  $d_k(u_1) \neq d_k(u_2)$  and  $v_1$  be any arbitrary vertex in  $G_2$ , then the number of vertices at distance  $k$  from  $(u_1, v_1)$  is given by

$$d_{k_{G_1 \oplus G_2}}(u_1, v_1) = d_k(u_1) \cdot d'_k(v_1) + d_k(u_1) \sum_{j=0}^{i-1} d'_j(v_1) + d'_k(v_1) \sum_{j=0}^{k-1} d_j(u_1)$$

and the number of vertices at distance  $k$  from  $(u_2, v_1)$  is given by

$$d_{k_{G_1 \oplus G_2}}(u_2, v_1) = d_k(u_2) d'_k(v_1) + d_k(u_2) \sum_{j=0}^{i-1} d'_j(v_1) + d'_k(v_1) \sum_{j=0}^{k-1} d_j(u_2),$$

implies  $d_{k_{G_1 \oplus G_2}}(u_1, v_1) \neq d_{k_{G_1 \oplus G_2}}(u_2, v_1)$ , since  $d_k(u_1) \neq d_k(u_2)$ . So  $G_1 \oplus G_2$  is not DDR, a contradiction. Hence  $G_1$  is DDR. Similarly we can prove  $G_2$  is also DDR. Hence, the proof.  $\square$

**Proposition 3.1:** *If the normal product  $G_1 \oplus G_2$  of two graphs is DDI then both  $G_1$  and  $G_2$  are DDI.*

*Proof.* Let the normal product  $G_1 \oplus G_2$  of two graphs be DDI. Suppose  $G_1$  is not DDI, then there exist two vertices  $u_1$  and  $u_2$  such that  $dds(u_1) = dds(u_2)$ . Let  $v_1$  be any vertex in  $G_2$ . The number of vertices at distance  $k, 0 \leq k \leq \max \{e(u_1), e(v_1)\}$  from  $(u_1, v_1)$  is given by

$$d_{k_{G_1 \oplus G_2}}(u_1, v_1) = d_k(u_1) d'_k(v_1) + d_k(u_1) \sum_{j=0}^{k-1} d'_j(v_1) + d'_k(v_1) \sum_{j=0}^{k-1} d_j(u_1).$$



and the number of vertices at distance  $k, 0 \leq k \leq \max\{e(u_2), e(v_1)\}$  from  $(u_2, v_1)$  is given by

$$d_{kG_1 \oplus G_2}(u_2, v_1) = d_k(u_2) d'_k(v_1) + d_k(u_2) \sum_{j=0}^{k-1} d'_j(v_1) + d'_k(v_1) \sum_{j=0}^{k-1} d_j(u_2).$$

Hence,  $d_{kG_1 \oplus G_2}(u_1, v_1) = d_{kG_1 \oplus G_2}(u_2, v_1), 0 \leq k \leq \max\{e(u_1) = e(u_2), e(v_1)\}$ , implies  $dds(u_1, v_1) = dds(u_2, v_1)$ , a contradiction. Hence  $G_1$  is DDI. Similarly we can prove  $G_2$  is also DDI. Hence, the proof.  $\square$

**Remark 3** : Normal product of two DDI graphs need not be DDI. The graphs in Figure 1 are the two DDI graphs whose normal product is not DDI.

**Lemma 3.2:** Let  $G_1$  and  $G_2$  be two DDI graphs. Let  $A = \{dds(u_i) | u_i \in G_1\}$  and  $B = \{dds(v_i) | v_i \in V(G_2)\}$ . If  $|A \cap B| \geq 2$  then  $G_1 \oplus G_2$  is not DDI.

*Proof.* Let  $|A \cap B| \geq 2$ . Then there exist  $u_k, u_l$  in  $G_1$  and  $v_m, v_n$  in  $G_2$  such that  $dds(u_k) = dds(v_m)$  and  $dds(u_l) = dds(v_n)$ . Hence in  $G_1 \oplus G_2$ ,  $dds(u_k, v_n) = dds(u_l, v_m)$ , making  $G_1 \oplus G_2$  non DDI. Hence the proof.  $\square$

**Proposition 3.2:** Let  $G_1$  and  $G_2$  be any two graphs. Let  $u$  be any vertex in  $G_1$  and  $S$  be a subset of  $V(G_2)$  such that no two vertices of  $S$  have same  $dds$ , then no two vertices of  $\{u\} \times S$  have same  $dds$  in  $G_1 \oplus G_2$ .

*Proof.* Suppose there exist two vertices  $(u, v)$  and  $(u, w)$  in  $\{u\} \times S$  having same  $dds$ , i.e.,  $dds(u, v) = dds(u, w)$ , this condition is satisfied only if  $e(v) = e(w)$ . Here three subcases arise, viz. **Case(a):**  $e(u) < e(v) = e(w)$ , **Case(b):**  $e(u) = e(v) = e(w)$  and **Case(c):**  $e(u) > e(v) = e(w)$ .

**Case(a):**  $dds(u, v) = dds(u, w), e(u) < e(v) = e(w), e(u, v) = e(v) = e(w)$ .

$$d_{mG_1 \oplus G_2}(u, v) = d_m(u) d'_m(v) + d_m(u) \sum_{j=0}^{m-1} d'_j(v) + d'_m(v) \sum_{j=0}^{m-1} d_j(u) \text{ and } \quad (1)$$

$$d_{mG_1 \oplus G_2}(u, w) = d_m(u) d'_m(w) + d_m(u) \sum_{j=0}^{m-1} d'_j(w) + d'_m(w) \sum_{j=0}^{m-1} d_j(u), \quad (2)$$

$d_{mG_1 \oplus G_2}(u, v) = d_{mG_1 \oplus G_2}(u, w)$ , for all  $m, 0 \leq m \leq e(v) = e(w)$ .

First taking,  $d_{mG_1 \oplus G_2}(u, v) = d_{mG_1 \oplus G_2}(u, w)$ , for all  $m, 0 \leq m \leq e(u)$ , we get

$$d_m(u)[d'_m(v) - d'_m(w)] + d_m(u) \left[ \sum_{j=0}^{m-1} d'_j(v) - \sum_{j=0}^{m-1} d'_j(w) \right] + \left[ \sum_{j=0}^{m-1} d_j(u) [d'_m(v) - d'_m(w)] \right] = 0$$

i.e.,  $\left[ d_m(u) + \sum_{j=0}^{m-1} d_j(u) \right] [d'_m(v) - d'_m(w)] + d_m(u) \left[ \sum_{j=0}^{m-1} d'_j(v) - \sum_{j=0}^{m-1} d'_j(w) \right] = 0$

i.e.,  $\left[ \sum_{j=0}^m d_j(u) \right] [d'_m(v) - d'_m(w)] + d_m(u) \left[ \sum_{j=0}^{m-1} d'_j(v) - \sum_{j=0}^{m-1} d'_j(w) \right] = 0 \quad (3)$

Put  $m = 1$  in Eqn (3).  $\implies [d_0(u) + d_1(u)][d'_1(v) - d'_1(w)] + d_1(u)[d'_0(v) - d'_0(w)] = 0$

i.e.,  $d'_1(v) = d'_1(w)$

Put  $m = 2$  in Eqn. (3).  $\implies \left[ \sum_{j=0}^2 d_j(u) \right] [d'_2(v) - d'_2(w)] + d_2(u)$

$$\left[ \sum_{j=0}^1 d'_j(v) - \sum_{j=0}^1 d'_j(w) \right] = 0$$

i.e.,  $d'_2(v) = d'_2(w)$ , And so on,

Put  $m = e(u)$  in Eqn. (1).  $\implies \left[ \sum_{j=0}^{e(u)} d_j(u) \right] [d'_{e(u)}(v) - d'_{e(u)}(w)]$

$$+ d_{e(u)}(u) \left[ \sum_{j=0}^{e(u)-1} d'_j(v) - \sum_{j=0}^{e(u)-1} d'_j(w) \right] = 0$$

i.e.,  $d'_{e(u)}(v) = d'_{e(u)}(w)$ .

Hence  $d'_m(v) = d'_m(w)$  for all  $m, 0 \leq m \leq e(u)$ . (4)

Now for all  $m, e(u) < m \leq e(v) = e(u, w)$ , we have

$$d_{m_{G_1 \oplus G_2}}(u, v) = d'_m(v) \sum_{j=0}^{e(u)} d_j(u) \text{ and } d_{m_{G_1 \oplus G_2}}(u, w) = d'_m(w) \sum_{j=0}^{e(u)} d_j(u).$$

Hence  $d_{m_{G_1 \oplus G_2}}(u, v) = d_{m_{G_1 \oplus G_2}}(u, w)$ , for all  $m, e(u) < m \leq e(v)$

$$= e(u, w) \text{ gives } d'_m(v) \sum_{j=0}^{e(u)} d_j(u) - d'_m(w) \sum_{j=0}^{e(u)} d_j(u) = 0, \text{ i.e.,}$$

$$[d'_m(v) - d'_m(w)] \sum_{j=0}^{e(u)} d_j(u) = 0 \text{ i.e., } d'_m(v) - d'_m(w) = 0.$$

$$\text{Hence } d'_m(v) = d'_m(w), \text{ for all } m, e(u) < m \leq e(v) = e(u, w). \quad (5)$$

Combining (4) and (5), we get  $d'_m(v) = d'_m(w)$ , for all  $m, 0 \leq m \leq e(v) = e(w) = e(u, w)$ , i.e.,  $dds(v) = dds(w)$ , a contradiction. Hence  $dds(u, v) \neq dds(u, w)$ .

**Case (b):**  $dds(u, v) = dds(u, w), e(u) = e(v) = e(w) = e(u, v)$ .

Substituting these conditions in (1) and (2), we get

$$\left[ \sum_{j=0}^m d_j(u) \right] [d'_m(v) - d'_m(w)] + d_m(u) \left[ \sum_{j=0}^{m-1} d'_j(v) - \sum_{j=0}^{m-1} d'_j(w) \right] = 0 \quad (6)$$

Substituting values of  $m$  in Eq.(6) we get  $d'_m(v) = d'_m(w)$  for all  $m, 0 \leq m \leq e(u) = e(v) = e(w) = e(u, v)$ .

i.e.,  $dds(v) = dds(w)$ , a contradiction. Hence  $dds(u, v) \neq dds(u, w)$ .

**Case (c):**  $e(u) > e(v) = e(w), e(u, v) = e(u)$ .

Substituting the values in (1) and (2) we get

$$\left[ \sum_{j=0}^m d_j(u) \right] [d'_m(v) - d'_m(w)] + d_m(u) \left[ \sum_{j=0}^{m-1} d'_j(v) - \sum_{j=0}^{m-1} d'_j(w) \right] = 0 \quad (7)$$

Substituting the values of  $m$  in (7) we get

$$d'_m(v) = d'_m(w) \text{ for all } m, 0 \leq m \leq e(v) = e(w).$$

i.e.,  $dds(v) = dds(w)$ , a contradiction. Hence,  $dds(u,v) \neq dds(u,w)$ .  $\square$

**Remark 4:** Let  $S_1$  and  $S_2$  be two subsets of  $V(G_1)$  such that every pair  $(x,y)$ ,  $x \in S_1, y \in S_2$  satisfies  $dds(x) \neq dds(y)$  and let  $z \in V(G_2)$  be any vertex in  $G_2$ , then in  $G_1 \oplus G_2$ , the subsets  $\{(x,z) \mid x \in S_1\}$  and  $\{(y,z) \mid y \in S_2\}$  are such that every pair  $((x,z), (y,z))$ ,  $x \in S_1$  and  $y \in S_2$  satisfies  $dds(x,z) \neq dds(y,z)$ .

#### 4. Acknowledgement

The third author (SAUS) thanks the UGC, New Delhi for granting fellowship under Special Assistance Program CAS Phase-II.

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*Received December, 2011*